



Multi-Layered Unstructured Mesh Generation

Panagiotis Foteinos^{1,2}
pfot@cs.wm.edu

Daming Feng²
dfeng@cs.odu.edu

Andrey Chernikov²
achernik@cs.odu.edu

Nikos Chrisochoides²
nikos@cs.odu.edu



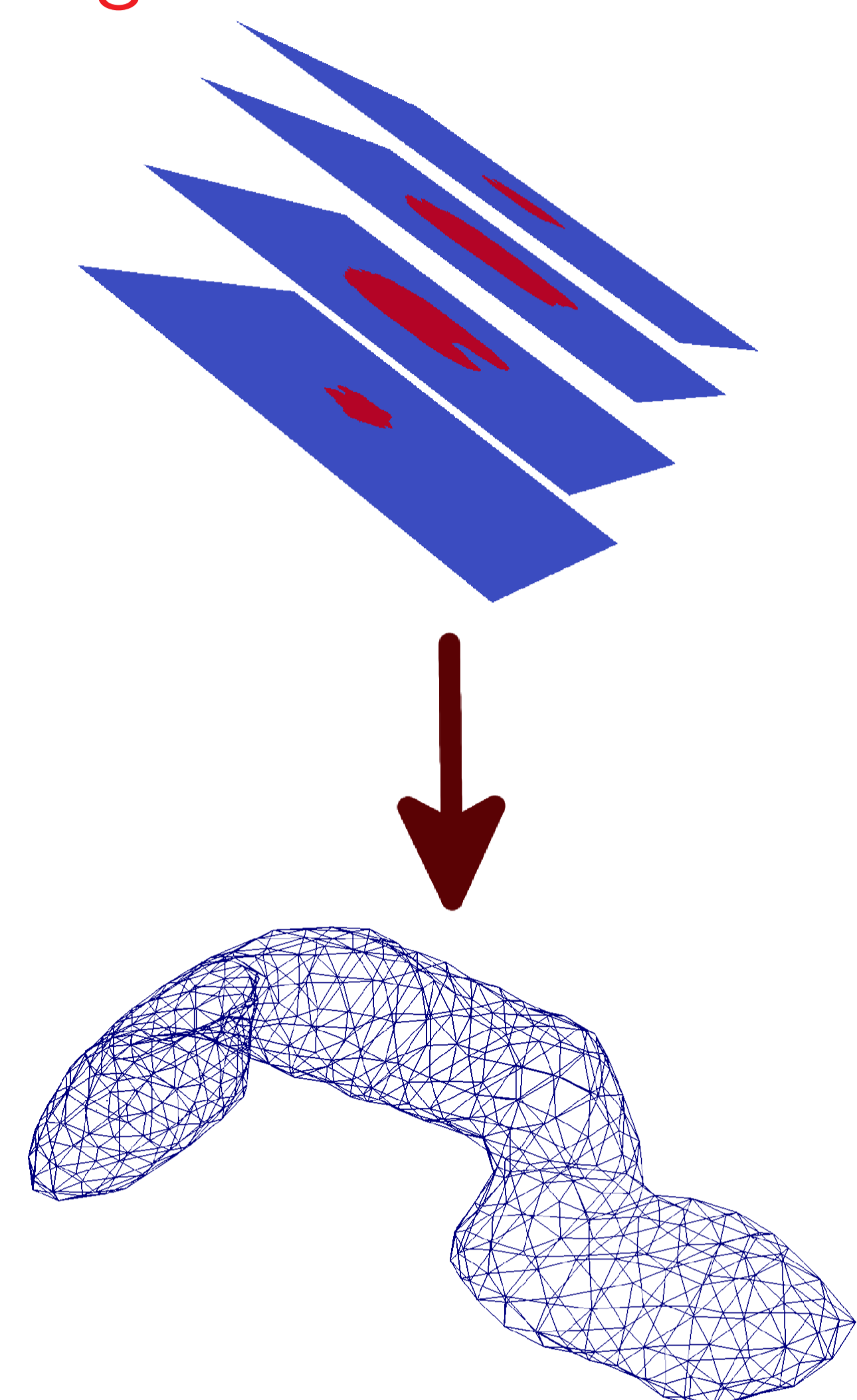
¹College of William and Mary, Computer Science, Williamsburg, VA, USA

²Old Dominion University, Computer Science, Norfolk, VA, USA

Introduction

- ▶ Input: **Image**
- ▶ Output: **Mesh**
- ▶ Faithful representation of the underlying object: **Fidelity**
- ▶ Well shaped tetrahedra: **Quality**
- ▶ Fidelity: symmetric Hausdorff distance, ambient isotopy
- ▶ Quality: aspect ratio, radius-edge ratio, size

Image to Mesh Conversion

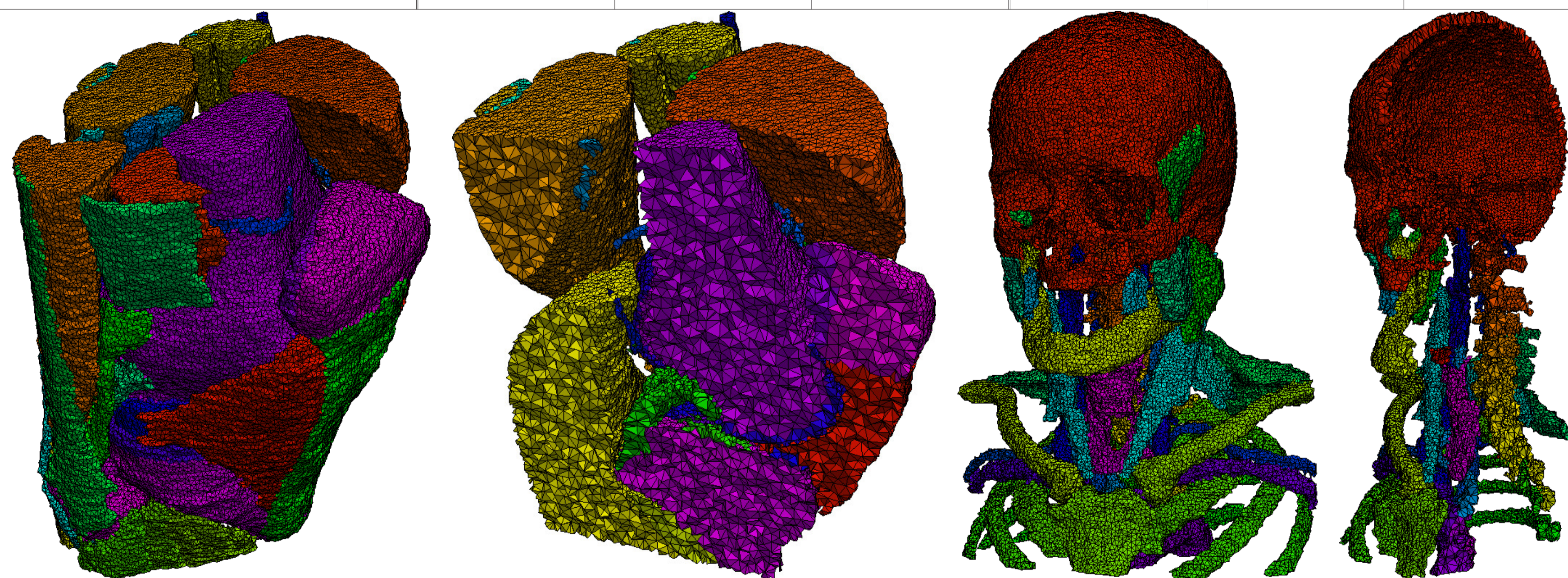


▶ **Goal: scalability on thousands of cores!**

Parallel Image-to-Mesh Conversion (PI2M)

Statistics regarding the single-threaded performance and the quality/fidelity achieved by PI2M and CGAL. PI2M includes the extra overhead introduced by synchronization, contention management, and load balancing to support the (potential) presence of other threads.

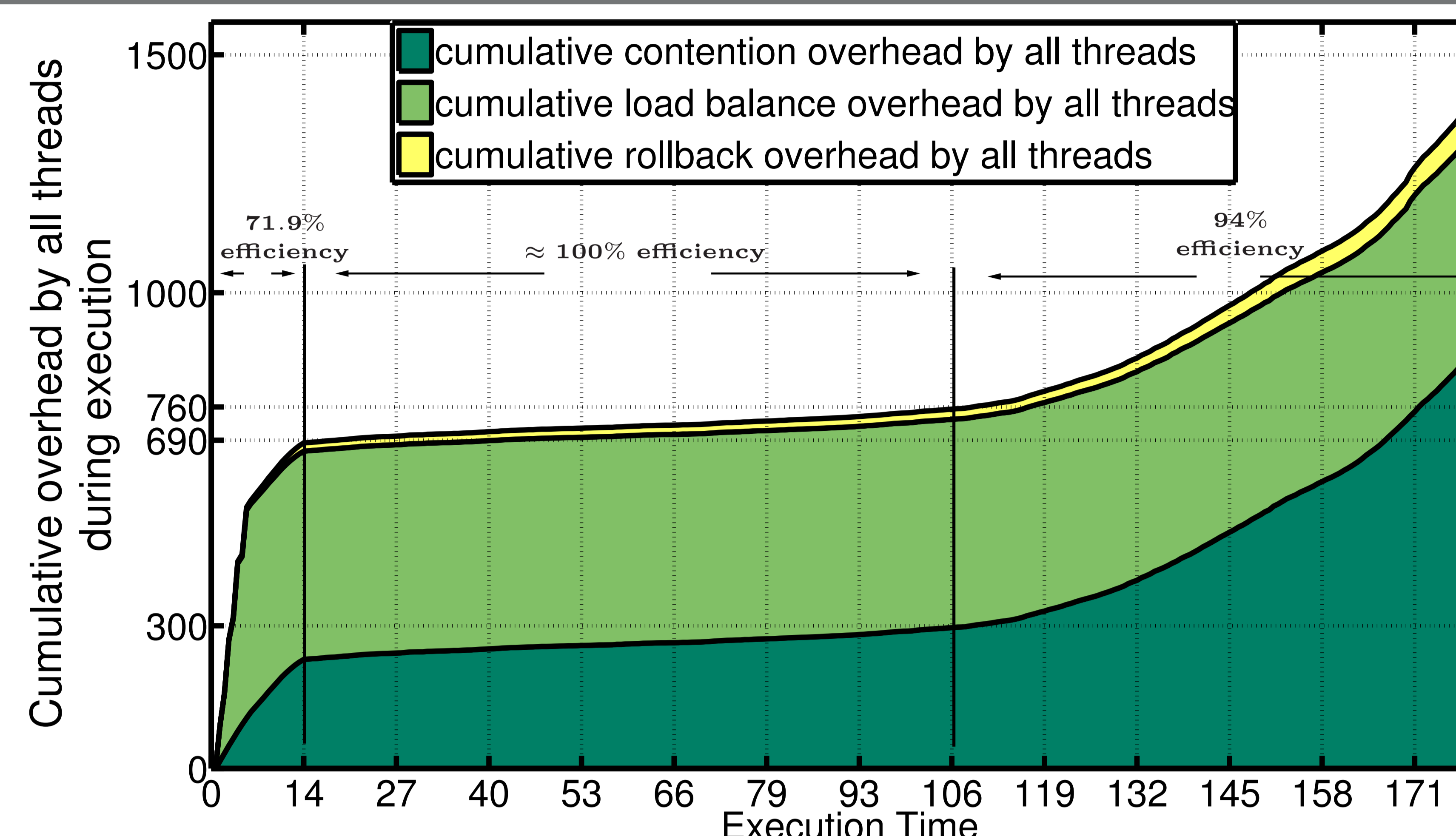
	knee atlas			head-neck atlas		
	PI2M	CGAL	TetGen	PI2M	CGAL	TetGen
#tetrahedra / seconds time	67,609 / 6.5 secs	40,069 / 10.9 secs	98,658 / 4.4 secs	96,464 / 10.3 secs	29,077 / 34.1 secs	61,903 / 16.0 secs
#tetrahedra	439,458	436,749	434,095	993,583	991,509	990,446
max radius-edge ratio	2	4.4	18.6	2	11.2	93.4
smallest boundary planar angle	17.4°	24.6°	18.0°	15.8°	2.4°	15.3°
(min, max) dihedral angles	(4.6°, 170.1°)	(2.5°, 176.3°)	(2.9°, 173.0°)	(4.5°, 170.2°)	(4.1°, 173.9°)	(0.4°, 172.0°)
Hausdorff distance	10.7 mm	10.3 mm	-	15.3 mm	15.2 mm	-



▶ **Excellent Single-Threaded performance**

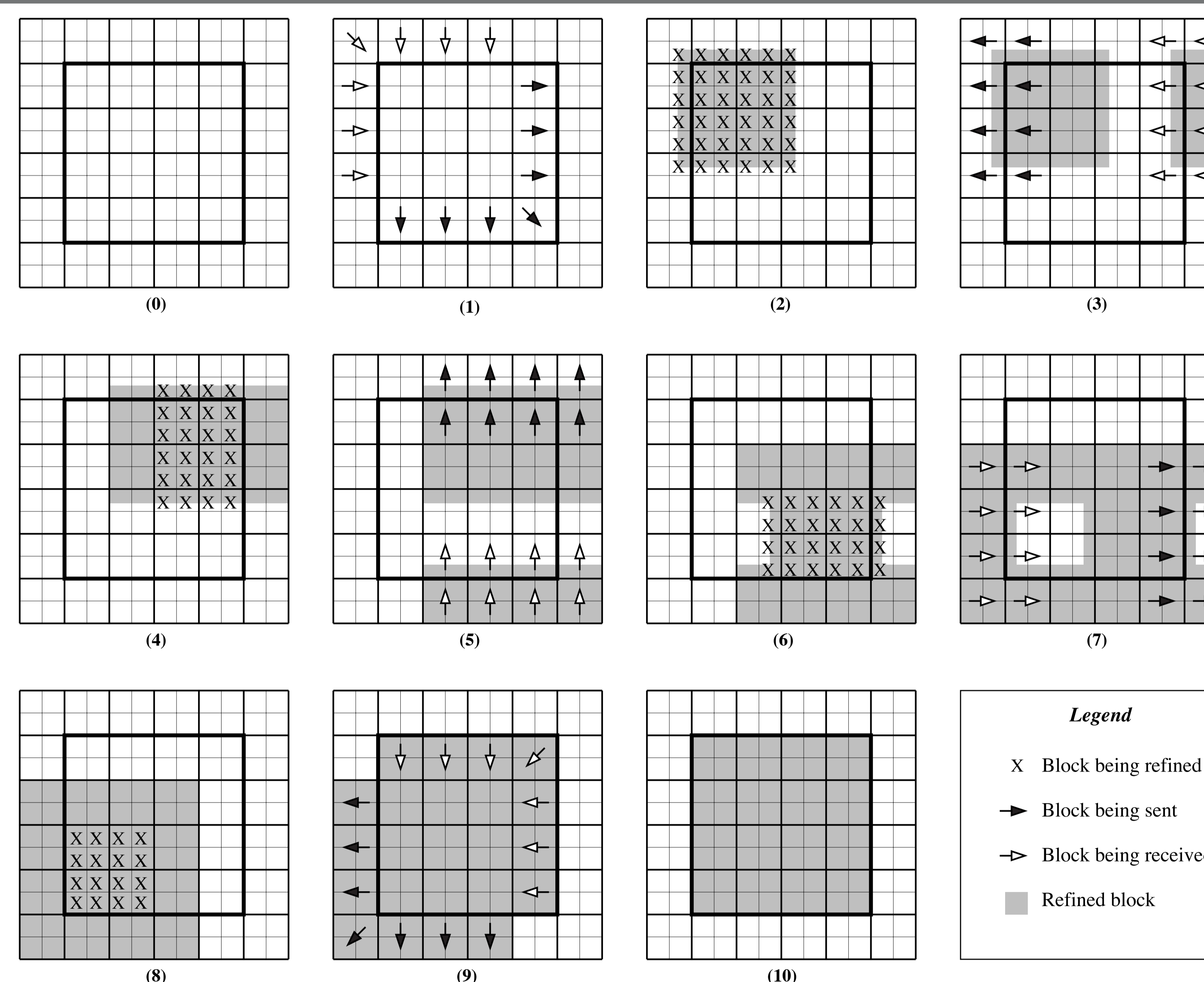
▶ **Scalability for up to 100 cores**

Bottleneck: memory latency



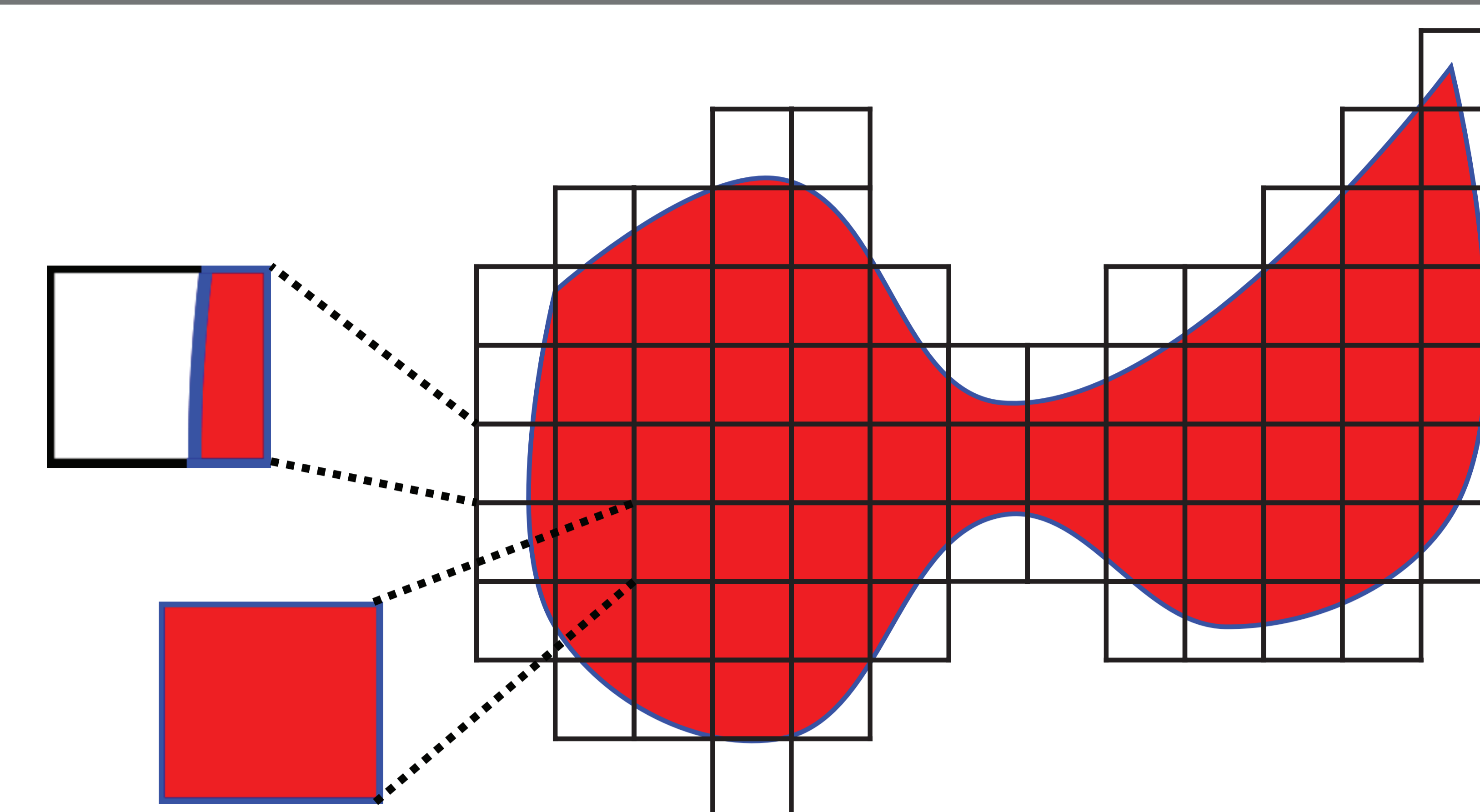
- ▶ Considering zero overhead for load balancing and contention, the 106s-14s=92s is far from perfect...
- ▶ Many small packages increase traffic pressure

Data Decomposition



- ▶ Impose an upper limit R for the elements' circumradii
- ▶ Data subdomains of size R
- ▶ Non-adjacent subdomains are safely independent
- ▶ Improves data locality
- ▶ 69% efficiency on 10 cores
- ▶ **Data Locality + PI2M: $10^2 \times 6.9 = 690$ concurrency in a rack of 10 nodes**

Domain Decomposition



- ▶ Data Decomposition alleviates intensive memory pressure, but it does not eliminate it
- ▶ Domain Decomposition separates memory banks
- ▶ Delaunay admissible medial axis domain decomposition is difficult in 3D or 4D
- ▶ Introduce artificial boundaries that do not hurt fidelity
- ▶ 66% efficiency on 48 cores
- ▶ **Domain Decomposition + Data Locality + PI2M: $10^2 \times 6.9 \times (0.66 \times 48) \approx 22,000$ concurrency in an enclosure of 48 racks**

Acknowledgements

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References

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